Security Weaknesses in Two Certificateless Signcryption Schemes

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Abstract. Recently, a certificateless signcryption scheme in the standard model was proposed by Liu et al. in [1]. Another certificateless signcryption scheme in the standard model was proposed by Xie et al. in [2]. Here, we show that the scheme in [1] and [2] are not secure against Type-I adversary.

1 Certificateless Signcryption Scheme by Liu et al.[1]

1.1 Review of the Scheme

In this section, we review the certificateless signcryption scheme secure against malicious-but-passive KGC attacks in the standard model proposed by Liu et al. The proposed scheme involves three parties: a KGC, a sender with an identity U_S and a receiver with an identity U_R. The scheme consists of the following algorithms.

Setup : Let $(\mathbb{G}, \mathbb{G}_T)$ be bilinear groups, where $|\mathbb{G}| = |\mathbb{G}_T = p$ for some prime p and g be a generator of $\mathbb{G}$. Let $\hat{e} : \mathbb{G} \times \mathbb{G} \rightarrow \mathbb{G}_T$ be the bilinear pairing and $H : \{0,1\}^* \rightarrow \mathbb{G}_T$ be the collision resistant hash function. KGC chooses randomly $\alpha \in \mathbb{Z}_p$ and computes $g_1 = g^\alpha$. Additionally, the KGC selects three random values $g_2, u', v' \in \mathbb{G}$ and two vectors $\mathcal{U} = (u_i)_n$, $\mathcal{V} = (v_j)_m$ whose elements are chosen from $\mathbb{G}$ at random. The system parameters are $params = (\mathbb{G}, \mathbb{G}_T, \hat{e}, g, g_1, g_2, u', v', \mathcal{U}, \mathcal{V}, H)$ and the master secret key is g_2^α.

Partial-Private-Key-Extract : Let $u[i]$ denote the i^{th} bit of an identity $u \in \{0,1\}^n$ and $\hat{u} = \{i | u[i] = 1, i = 1, \ldots, n\}$. The KGC picks $r \in \mathbb{Z}_p$ uniformly and computes,

$$d_u = (d_{u,1}, d_{u,2}) = (g_2^\alpha \left(u' \prod u_i\right)^r, g^r).$$

An entity with identity u is given d_u as his partial private key. Therefore, the sender and the receivers partial private keys are,

$$d_S = (d_{S,1}, d_{S,2}) = (g_2^\alpha \left(u' \prod u_i\right)^{r_S}, g^{r_S}).$$

$$d_R = (d_{R,1}, d_{R,2}) = (g_2^\alpha \left(u' \prod u_i\right)^{r_R}, g^{r_R}).$$

User-Key-Generate : An entity with an identity u chooses randomly a secret value $x_u \in \mathcal{Z}_p$ and computes a public key,

$$pk_u = \hat{e}(g_1, g_2)^{x_u}$$

Private-Key-Extract : An entity with identity u picks $r' \in \mathbb{Z}_p$ at random, and computes a private key,

$$sk_u = sk_{u,1}, sk_{u,2} = \left(d_{u,1}^{x_u} \left(u' \prod u_i\right)^{r'}, d_{u,2}^{x_u} g^{r'} \right)$$

where $t = rx_u + r'$.

Signcrypt : To send a message $M \in \mathbb{G}_T$ to the receiver with public key $pk_R = \hat{e}(g_1, g_2)^{x_R}$, the sender picks $r'' \in \mathbb{Z}_p$ randomly and carries out the following steps.

- Compute $\sigma_1 = M.pk_R^{r''} = m.\hat{e}(g_1, g_2)^{x_R r''}$.
- Compute $\sigma_2 = g^{r''}$.
- Compute $\sigma_3 = (u' \prod u_i)^{r''}$.
- Set $\sigma_4 = sk_{S,2}$
- Compute $\hat{M} = H(\sigma_1, \sigma_2, \sigma_3, \sigma_4, u_R, pk_R) \in \{0,1\}^m$, where $m[j]$ denotes the j^{th} bit of $\hat{M}$ and $\mathcal{M} = \{j | m[j] = 1, j = 1, 2, \ldots, m\}$.
- Compute $\sigma_5 = sk_{S,1}.(v' \prod v_i)^{r''}$.
- Output the ciphertext $\sigma = (\sigma_1, \sigma_2, \sigma_3, \sigma_4, \sigma_5)$.

Unsigncrypt : Upon receiving a ciphertext $\sigma = (\sigma_1, \sigma_2, \sigma_3, \sigma_4, \sigma_5)$, the receiver decrypts the ciphertext as follows.

- Compute $\hat{M} = H(\sigma_1, \sigma_2, \sigma_3, \sigma_4, u_R, pk_R) \in \{0,1\}^m$, where $m[j]$ denotes the j^{th} bit of $\hat{M}$ and $\mathcal{M} = \{j | m[j] = 1, j = 1, 2, \ldots, m\}$.
- Check that the equality,

$$\hat{e}(\sigma_5, g) = pk_S.\hat{e}(u' \prod u_i, \sigma_4)\hat{e}(v' \prod v_j, \sigma_2)$$

holds. If not output "$Invalid$". Otherwise, compute and output $M = \sigma_1.\hat{e}(\sigma_3, sk_{R,2})/\hat{e}(\sigma_2, sk_{R,1})$

1.2 Attack on the Scheme by Liu et al. :

The scheme proposed by Liu et al. in [1] does not provide confidentiality against Type-I adversary. We show the scheme is not even CPA secure against Type-I adversary. The attack can be launched by a Type-I adversary by replacing the public key of the target receiver whose signcryption the adversary wants to designcrypt. This can be achieved in the following way :
During the Type-I confidentiality game,

- The challenger runs the setup and provides the system public parameters to the adversary.
- The adversary has access to all the oracles namely **Partial-Private-Key-Extract**, **Private-Key-Extract**, **Replace-Public-Key**, **Signcrypt** and **Unsigncrypt**.
- The adversary replaces the public key of the receiver (say R^*) which he wants to use during the *challengephase* by $pk_{R^*} = \hat{e}(g, g)^{r^*}$ where $r^* \in_R \mathbb{Z}_p$.
- Without asking any further queries the adversary now picks two messages $\{m_0, m_1\}$ of equal length and a sender identity S and receiver identity R^* on which the adversary wishes to be challenged and sends to the challenger.

– The Challenger now picks a random bit $\delta \in \{0,1\}$, cooks up the signcryption $\sigma^* = (\sigma_1^*, \sigma_2^*, \sigma_3^*, \sigma_4^*, \sigma_5^*)$ of message m_δ and sends σ^* to the adversary.
– Now the adversary can get back the key by performing $m_{\delta'} = \sigma_1^* \hat{e}(\sigma_2, g^{r^*})$ and outputs δ' to the challenger.
– Hence the adversary can successfully distinguish the message being signcrypted. This clearly shows that the scheme given by Liu et al. is not CPA secure against Type-I adversary.

2 Certificateless Signcryption Scheme by Xie et al.[2]

Since the scheme is available in public medium, we do not review the scheme here.

Attack on the Scheme

In this section we present a total break of the certificateless signcryption scheme in [2] by Type-I adversary. During the unforgeability game, the adversary knows the full private key of the receiver. Thus, during the training phase, the Type-I forger queries and obtains a ciphertext $\sigma = \langle c, u, v, w \rangle$ from the signcrypt oracle. Let σ be a signcryption from sender ID_A to receiver ID_B, where the private key D_B corresponding to the receiver is known to the adversary. The adversary performs the following to compute the partial private key d_A of the sender.

– We know that $w = x_A h_2 + r_1$. (It is known that Type-I adversary can replace the public key and hence have access to the sender secret value x_A.)
– Computes $g^{r_1'} = \hat{e}(d_B, u)$ and $m = c \oplus H_3(g^{r_1'}, x_B u)$.
– Computes $h_2 = H_2(m, u, g^{r_1'}, x_B u, pk_A, pk_B)$.
– Computes $r_1 = w - x_A h_2$.
– It is now possible to compute $d_A = v\left(\frac{r_1 - h_2}{r_1}\right)$.

Hence, a Type-I adversary can find out the partial private key of any legitimate user in the system, which leads to a total break of the system in [2].

References

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